

# SNU 4541.664A Program Analysis Spring 2005 Note 5

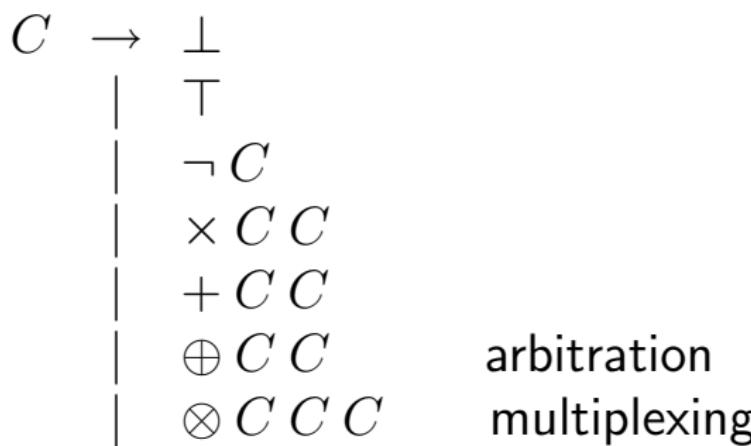
Prof. Kwangkeun Yi

# 계획

1 의미구조 정의하기 연습

2 프로그램의 의미는 아직도 고정점

# 회로도



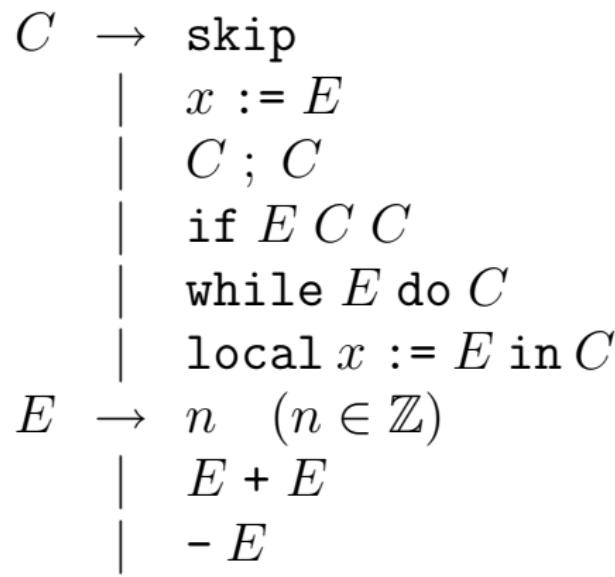
- by structural operational semantics
- by evaluation context semantics
- by transition semantics

C----

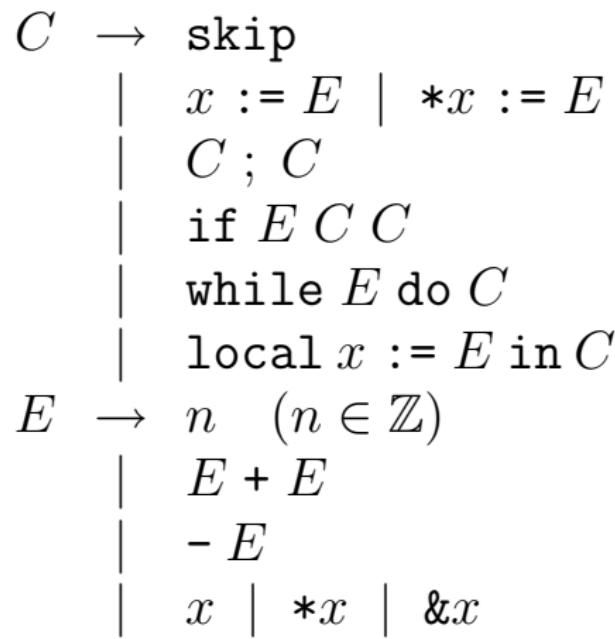
$$\begin{array}{lcl} C & \rightarrow & \text{skip} \\ & | & x := E \\ & | & C ; C \\ & | & \text{if } E \ C \ C \\ & | & \text{while } E \text{ do } C \\ E & \rightarrow & n \quad (n \in \mathbb{Z}) \\ & | & E + E \\ & | & - E \end{array}$$

- by evaluation context semantics
- by abstract machine semantics

C---



C--



C---+

$C$	$\rightarrow$	skip
		$x := E$
		$C ; C$
		$\text{if } E \ C \ C$
		$\text{while } E \text{ do } C$
		$\text{raise}$
		$\text{try } C \text{ handle } C$
$E$	$\rightarrow$	$n \quad (n \in \mathbb{Z})$
		$E + E$
		$- E$

# 의미는 아직도 고정점

- 한 실행과정:  $M \vdash C \Rightarrow M'$ 의 증명 또는  
 $(M, C) \rightarrow (M_1, C_1) \rightarrow \dots$
- $\llbracket C \rrbracket M = fix\lambda t.(M, C) \sqcup (t \rightarrow (M_{i+1}, C_{i+1}))$   
where  $t = \dots \rightarrow (M_i, C_i)$   
 $\wedge (M_i, C_i) \rightarrow (M_{i+1}, C_{i+1})$
- $\llbracket C \rrbracket = fix\lambda f.\lambda M.(M, C) \sqcup (f(M) \rightarrow (M_{i+1}, C_{i+1}))$   
where  $f(M) = \dots \rightarrow (M_i, C_i)$   
 $\wedge (M_i, C_i) \rightarrow (M_{i+1}, C_{i+1})$

$$\llbracket C \rrbracket M = fix \lambda T.$$

$$\{(M, C)\} \sqcup$$

$$\{t \rightarrow (M_{i+1}, C_{i+1}) \mid$$

$$t \in T, t = \dots \rightarrow (M_i, C_i), (M_i, C_i) \rightarrow (M_{i+1}, C_{i+1})\}$$

$$\llbracket C \rrbracket = fix \dots$$

$$\llbracket C \rrbracket M = fix \lambda S.$$

$$\{(M, C)\} \sqcup$$

$$\{(M_{i+1}, C_{i+1}) \mid$$

$$(M_i, C_i) \rightarrow (M_{i+1}, C_{i+1}), (M_i, C_i) \in S\}$$

$$\llbracket C \rrbracket = fix \dots$$